# Algorithms and Networking for Computer Games

Chapter 3: Tournaments

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#### Tournament types

- rank adjustment (or challege) tournament
  - each match is a challenge for a rank exchange
  - types: ladder, hill climbing, pyramid, king of the hill
- elimination tournament (or cup)
  - each match eliminates the loser from the tournament
  - types: random selection, random pairing, single elimination
- scoring tournament
  - each match rewards the winner
  - types: round robin
- hybridizations

#### Other uses for tournaments

#### game balancing

- duelling synthetic players
- adjusting point rewarding schemes

#### heuristic search

- selecting suboptimal candidates for a genetic algorithm
- group behaviour
  - modelling pecking order
- learning player characteristics
  - managing history knowledge



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### Example: Scoring tournament

	Tuomas	Ааро	Simeoni	Timo	Lauri	Eero
Juhani	m <sub>0</sub>	m <sub>c</sub>	<i>m</i> <sub>11</sub>	m <sub>15</sub>	m <sub>18</sub>	m <sub>20</sub>
	Tuomas	m <sub>1</sub>	m <sub>7</sub>	m <sub>12</sub>	m <sub>16</sub>	<i>m</i> <sub>19</sub>
		Ааро	<i>m</i> 2	111 <sub>8</sub>	m <sub>13</sub>	<i>m</i> <sub>17</sub>
			Simeoni	m <sub>3</sub>	m <sub>9</sub>	m <sub>14</sub>
				Timo	m <sub>4</sub>	<i>m</i> <sub>10</sub>
					Lauri	<i>m</i> 5

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#### Terms

- players:  $p_0 \dots p_{n-1}$
- match between  $p_i$  and  $p_j$ : match(i, j)
- outcome: win, lose, TIE
- rank of  $p_i$ : rank(i)
- players with the rank r: rankeds(r)
- round: a set of (possibly) concurrent matches
- bracket: diagram of match pairings and rounds

#### Rank adjustment tournaments

- a set of already ranked players
- matches
  - independent from one another
  - outcome affects only the participating players
- suits on-going tournaments
  - example: boxing
- matches can be limited by the rank difference

# Ladder and pyramid tournaments $p_i: rank(i) = 0$ $p_i: rank(j) = 1$ $p_i: rank(j) = 1$

 $p_k$ : rank(k) = 2 -

 $p_m$ : rank(m) = 2  $p_n: rank(n) = 2$ 

 $rankeds(2) = \{k, m, n\}$ 

 $p_k$ : rank(k) = 2

 $p_n: rank(n) = 4$ 

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#### **Elimination tournaments**

- loser of a match is eliminated from the tournament
  - $\blacksquare$  no ties!  $\rightarrow$  tiebreak competition
- winner of a match continues to the next round
- how to assign pairings for the first round?
  - seeding
- examples
  - football cups, snooker tournaments

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# Random selection and random pairing



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### Single elimination



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- some match pairing will not occur in a single elimination tournament
- pairings for the first round (i.e., seeding) affects the future pairings
- seeding can be based on existing ranking
  - favour the top-ranked players
  - reachability: give the best players an equal opportunity to proceed the final rounds

# Seeding methods

#### ■ random

- does not favour any player
- does not fulfil reachability criterion
- standard and ordered standard
  - favours the top-ranked players
  - ordered standard: matches are listed in increasing order

#### equitable

 in the first round, the rank difference between the players is the same for each match

#### Scoring tournaments

- round robin: everybody meets everybody else once
- scoring table determines the tournament winner
  players are rewards with scoring points
  winner and tie
- matches are independent from one another

#### Reduction to a graph

- *n* players
- clique  $K_n$
- players as vertices, matches as edges
- how to organize the rounds?
  - a player has at most one match in a round
  - a round has as many matches as possible



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## Reduction to a graph (cont'd)

- if *n* is odd, partition the edges of the clique to (n-1)/2 disjoint sets
  - in each turn, one player is resting
  - Implayer  $p_i$  rests in the round *i*
- if *n* is even, reduce the problem
  - player  $p_{n-1}$  is taken out from the clique
  - solve the pairings for n 1 players as above
  - for each round, pair the resting player  $p_i$  with player  $p_{n-1}$

#### Round robin with seven players

round		resting		
0	1 – 6	2-5	3 – 4	0
1	2-0	3 – 6	4 – 5	1
2	3 – 1	4-0	5 – 6	2
3	4 – 2	5 – 1	6 – 0	3
4	5 – 3	6 – 2	0 – 1	4
5	6 – 4	0-3	1 – 2	5
6	0-5	1 – 4	2-3	6

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#### Normalized round robin

- who is the resting player in a given round?
  → answered
- given two players, in which round they will face one another?
  - $\rightarrow$  no simple rule?
- change the selection of the resting player
  resting player: r · [(n + 1) / 2] mod n
  if n is odd, p<sub>i</sub> and p<sub>j</sub> will face in the round i + j mod (number of rounds)

### Real-world tournament examples

boxing

reigning champion and challengers

sport wrestling

double elimination: consolation bracket

professional wrestling

royal rumble

World Cup

ice hockey championship

■ snooker

#### **Practical considerations**

- home matches
- venue bookings
- travelling times
- risk management
- other costs