

## Cost function $g$

- actual cost from  $s$  to  $v$  along the cheapest path found so far
  - exact cost if  $G$  is a tree
  - can never underestimate the cost if  $G$  is a general graph
- $f(v) = g(s, v)$  and unit cost  $\rightarrow$  breadth-first search
- $f(v) = -g(s, v)$  and unit cost  $\rightarrow$  depth-first search

## Heuristic function $h$

- carries information from outside the graph
- defined for the problem domain
- the closer to the actual cost, the less superfluous vertices are expanded
- $f(v) = g(s, v) \rightarrow$  cheapest-first search
- $f(v) = h(v, r) \rightarrow$  best-first search

## Admissibility

- let Algorithm A be a best-first search using the evaluation function  $f$
- search algorithm is *admissible* if it finds the minimal path (if it exists)
  - if  $f = f^*$ , Algorithm A is admissible
- Algorithm A\* = Algorithm A using an estimate function  $h$ 
  - A\* is admissible, if  $h$  does not overestimate the actual cost

## Monotonicity

- $h$  is locally admissible  $\rightarrow h$  is monotonic
- monotonic heuristic is also admissible
- actual cost is never less than the heuristic cost  $\rightarrow f$  will never decrease
- monotonicity  $\rightarrow$  A\* finds the shortest path to any vertex the first time it is expanded
  - if a vertex is rediscovered, path will not be shorter
  - simplifies implementation

## Optimality

- Optimality theorem: The first path from  $s$  to  $r$  found by A\* is optimal.
- Proof: lecture notes p. 49

## Informedness

- the more closely  $h$  approximates  $h^*$ , the better A\* performs
- if  $A_1$  using  $h_1$  will never expand a vertex that is not also expanded by  $A_2$  using  $h_2$ ,  $A_1$  is more informed than  $A_2$
- informedness  $\rightarrow$  no other search strategy with *the same amount of outside knowledge* can do less work than A\* and be sure of finding the optimal solution

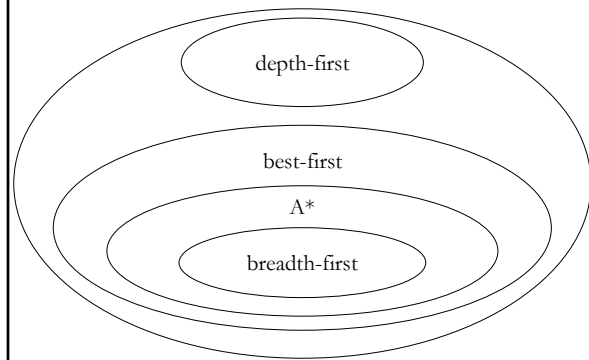
## Algorithm A\*

- because of monotonicity
  - all weights must be positive
  - closed list can be omitted
- the path is constructed from the mapping  $\pi$  starting from the goal vertex
  - $s \rightarrow \dots \rightarrow \pi(\pi(r)) \rightarrow \pi(r) \rightarrow r$

## Practical considerations

- computing  $h$ 
  - despite the extra vertices expanded, less informed  $h$  may yield computationally less intensive implementation
- suboptimal solutions
  - by allowing overestimation A\* becomes inadmissible, but the results may be good enough for practical purposes

## Search algorithms



## Realizing the movement

- movement through the waypoints
  - unrealistic: does not follow the game world geometry
  - aesthetically displeasing: straight lines and sharp turns
- improvements
  - hierarchical pathing
  - line-of-sight testing
- combining path finding to user-interface
  - real-time response

## Recapitulation

- discretization of the game world
  - grid, navigation mesh
  - waypoints, connections, costs
- path finding in a graph
  - Algorithm A\*
- realizing the movement
  - geometric corrections
  - aesthetic improvements

## Question

- Although this is the *de facto* approach in (commercial) computer games, are there alternatives?
- possible answers
  - AI processors (unrealistic?)
  - robotics: reactive agents (unintelligent?)
  - analytical approaches (inaccessible?)