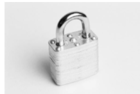
§10 Cheating Prevention

- traditional cheating in computer games
  - cracking the copy protection
  - fiddling with the binaries: boosters, trainers, etc.
- here, the focus is on multiplayer online games
  - exploiting technical advantages
  - exploiting social advantages
- cheaters' motivations
  - vandalism and dominance
  - peer prestige
  - greed

The goals of cheating prevention



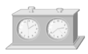
- protect the sensitive information
  - cracking passwords
  - pretending to be an administrator
- provide a fair playing field
  - tampering the network traffic
  - colluding with other players
- uphold a sense of justice inside the game world
  - abusing beginners
  - gangs

Network Security



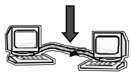
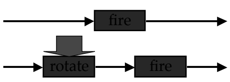
- Military
  - private networks → no problem
- Business, industry, e-commerce,…
  - 'traditional' security problems
- Entertainment industry
  - multiplayer computer games, online games
  - specialized problems

Taxonomy of Online Cheating 1 (4)

- Cheating by compromising passwords
  - dictionary attacks
- Cheating by social engineering
  - password scammers
- Cheating by denying service from peer players
  - denial-of-service (DoS) attack
  - clog the opponent's network connection

Taxonomy of Online Cheating 2 (4)

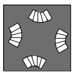



- Cheating by tampering with the network traffic
  - reflex augmentation
  - packet interception
  - look-ahead cheating
  - packet replay attack
- Cheating with authoritative clients
  - receivers accept commands blindly
    - requests instead of commands
    - checksums from the game state

Taxonomy of Online Cheating 3 (4)


- Cheating due to illicit information
  - access to replicated, hidden game data
  - compromised software or data
- Cheating related with internal misuse
  - privileges of system administrators
  - logging critical operations into CD-ROMs
- Cheating by exploiting a bug or design flaw
  - repair the observed defects with patches
  - limit the original functionality to avoid the defects
  - good software design in the first place!

Taxonomy of Online Cheating 4 (4)



- Cheating by collusion
  - two or more players play together without informing the other participants
  - one cheater participates as two or more players

- Cheating related to virtual assets
  - demand ⇒ supply ⇒ market ⇒ money flow ⇒ cheating
- Cheating by offending other players
  - acting against the 'spirit' of the game

Breaking the control protocol: Maladies & remedies

- *malady*: change data in the messages and observe effects
- *remedy*: checksums (MD5 algorithm)
- *malady*: reverse engineer the checksum algorithm
- *remedy*: encrypt the messages
- *malady*: attack with packet replay
- *remedy*: add state information (pseudo-random numbers)
- *malady*: analyse messages based on their sizes
- *remedy*: modify messages and add a variable amount of junk data to messages

MD5 algorithm

- message digest = a constant length 'fingerprint' of the message
- no one should be able to produce
  - two messages having the same message digest
  - the original message from a given message digest
- R. L. Rivest: MD5 algorithm
  - produces a 128-bit message digest from an arbitrary length message
- collision attack: different messages with the same fingerprint
- finding collisions is (now even technically!) possible
  - what is the future of message digest algorithms?

Illicit information

- access to replicated, hidden game data
  - removing the fog of war
  - compromised graphics rendering drivers
- cheaters have more knowledge than they should have → passive cheating
- compromised software or data
- counter-measures in a networked environment
  - centralized: server maintains integrity among the clients
  - distributed: nodes check the validity of each other's commands to detect cheaters

TOP SECRET

Exploiting design defects

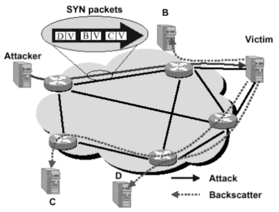
- what can we do to poor designs!
  - repair the observed defects with patches
  - limit the original functionality to avoid the defects
- client authority abuse
  - information from the clients is taken face-value regardless its reliability
- unrecognized (or unheeded) features of the network
  - operation when the latencies are high
  - coping with DoS and other attacks

Denial-of-Service (DoS) Attack

- Attack types:
  - logic attack: exploit flaws in the software
  - flooding attack: overwhelm the victim's resources by sending a large number of spurious requests
- Distributed DoS attack: attack simultaneously from multiple (possibly cracked) hosts
- IP spoofing: forge the source address of the outgoing packets
- Consequences:
  - wasted bandwidth, connection blockages
  - computational strain on the hosts

Analysing DoS Activity

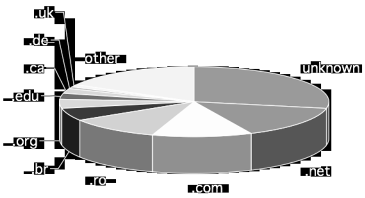
- Backscatter analysis
- Spoofing using random source address
- A host on the Internet receives unsolicited responses
- An attack of m packets, monitor n addresses
- Expectation of observing an attack: $E(X) = nm/2^{32}$



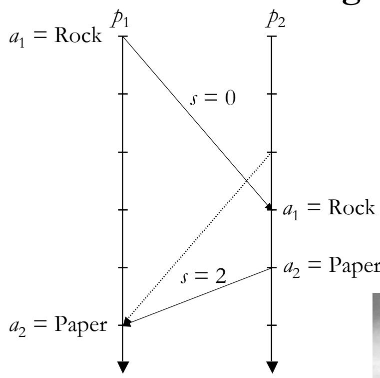
DoS: Selected Results

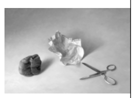
- Three week-long logging periods, February 2001
- >12,000 attacks, >5,000 distinct targets
- Significant number of attacks were directed against
  - home machines
  - users running Internet Relay Chat (IRC)
  - users with names that are sexually suggestive or incorporate themes of drug use
  - users supporting multiplayer games
- In addition to well-known Internet sites, a large range of small and medium sized businesses were targeted

DoS: Most Attacked Top-Level Domains



Look-ahead cheating





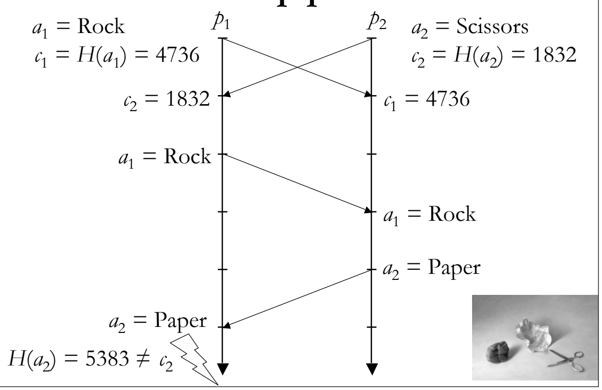
Two problems

- delaying one's decision
  - announce own action only after learning the opponent's decision
  - one-to-one and one-to-many
- inconsistent decisions
  - announce different actions for the same turn to different opponents
  - one-to-many

Lockstep protocol

1. Announce a commitment to an action.
   - commitment can be easily calculated from the action but the action cannot be inferred from the commitment
   - formed with a one-way function (e.g., hash)
2. When everybody has announced their commitments for the turn, announce the action.
   - everybody knows what everybody else has promised to do
3. Verify that the actions correspond to the commitments.
   - if not, then somebody is cheating…

Lockstep protocol



Loosening the synchronization 1(2)

- the slowest player dictates the speed
  - short turns
  - time limits for the announcements
- asynchronous lockstep protocol
  - sphere of influence: synchronization is needed only when the players can affect each other in the next turn(s)
  - otherwise, the players can proceed asynchronously

Loosening the synchronization 2(2)

- pipelined lockstep protocol
  - player can send several commitments which are pipelined
  - drawback: look-ahead cheating if a player announces action earlier than required
- adaptive pipeline protocol
  - measure the actual latencies between the players
  - grow or shrink the pipeline size accordingly

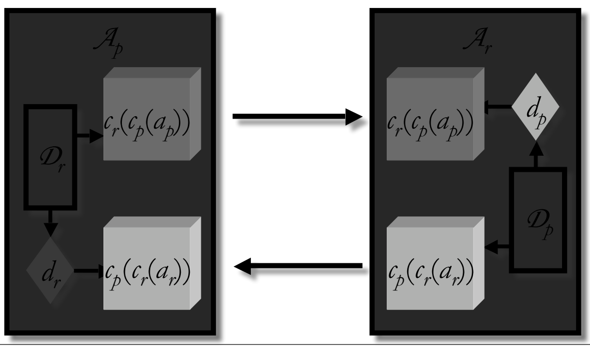
Drawbacks of the lockstep protocol

- requires two separate message transmissions
  - commitment and action are sent separately
  - slows down the communication
- requires a synchronization step
  - the slowest player dictates the pace
    - improvements: asynchronous lockstep, pipelined lockstep, adaptive pipeline lockstep
- does not solve the inconsistency problem!

Idea #1: Let's get rid of the repeat!

- send only a single message
  - but how can we be sure that the opponent cannot learn the action before annoucing its own action?
- the message is an active object, a *delegate*
  - program code to be run by the receiver (host)
  - delegate's behaviour cannot be worked out by analytical methods alone
  - guarantees the message exchange on a possibly hostile environment
- delegate provides the action once the host has sent its own action *using* the delegate

Example with two players

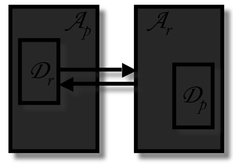


Threats

- what if the host delays or prevents the delegate's message from getting to its originator?
  - the host will not receive the next delegate until the message is sent
- what if the originator is malicious and the delegate spies or wastes the host's resources?
  - sandbox: the host restricts the resources available to the delegate
- how can the delegate be sure that it is sending messages to its originator?
  - communication check-up

Communication check-up

- the delegate sends a unique identification to its originator
  - static and dynamic information
- the delegate waits until the originator has responded correctly
- check-ups are done randomly
  - probability can be quite low
  - host cannot know whether the transmission is the actual message or just a check-up



Idea #2: Peer pressure

- players gossip the other players' actions from the previous turn(s)
- compare gossip and recorded actions; if there are inconsistencies, ban the player
  - cheating is detected only afterwards
  - gossiping imposes a threat of getting caught
- gossip is piggybacked in the ordinary messages
  - no extra transmissions are required
- how to be sure that the gossip is not forged?
  - rechecking with randomly selected players

How much is enough?

- example: 10 players, 60 turns, 1 cheater who forges 10% of messages, gossip from one previous turn
  - 1% gossip: *P*(cheater gets caught) = 0.44
  - 5% gossip: *P*(cheater gets caught) = 0.91
  - 10% gossip: *P*(cheater gets caught) = 0.98
- example: 100 players, 60 turns, 1 cheater who forges 10% of messages
  - 1% gossip: *P*(cheater gets caught) = 0.98
- example: 10 players, 360 turns, 1 cheater who forges 10% of messages
  - 1% gossip: *P*(cheater gets caught) = 0.97

Message

- action for the current turn t
- delegate for the next turn $t + 1$
- set of actions (i.e., gossip) from the previous turn $t - 1$